

## HOMEWORK #2

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**(5 points)** Motivation: In functional languages, we can think of data structures as directed acyclic graphs (DAGs), where each node (representing a struct) can contain only  $O(1)$  data and  $O(1)$  pointers to other nodes. We cannot change any nodes, we can only create new nodes. How can we implement a double ended queue a.k.a. deque in such languages?

One way is to represent the deque using two stacks (which are easy to implement as linked lists). For example, a double ended queue with elements

1, 2, 3, 4, 5, 6, 7, 8, 9

can be represented using 2 stacks:

**front** :     $\vdash$  3, 2, 1  
**rear** :     $\vdash$  4, 5, 6, 7, 8, 9

( $\vdash$  is the bottom of the stack, the top is on the right).

If we want to add/remove an element from the beginning of the queue, we add/remove it from **front**, if we want to add/remove an element from the end, we add/remove it from **rear**. The only problem is when one of the stacks gets empty. For example, if we remove elements 1, 2, 3 from the beginning, we get an empty **front**:

**front** :     $\vdash$   
**rear** :     $\vdash$  4, 5, 6, 7, 8, 9

What if we now want to take the element from the beginning again?

- a) One idea is to take all elements from **rear** and insert them in reverse order into **front**. We solve the removal and insertion at the other end symmetrically. Show that this is not a good implementation and prove that it *doesn't* work in  $O(1)$  *amortized time*. What is the lower bound?
- b) Better idea<sup>1</sup> is to maintain the invariant that
  - $size(\mathbf{front}) \leq 4size(\mathbf{rear}) + 1$  and
  - $size(\mathbf{rear}) \leq 4size(\mathbf{front}) + 1$ .

If the invariant is violated during any operation, we rearrange the stacks so that both stacks have the same size ( $\pm 1$  if the number of elements is odd). Prove that this algorithm runs in  $O(1)$  *amortized time for each operation*.

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<sup>1</sup>see e.g. `Data.Deque` library in Haskell, which works exactly like this:  
<https://hackage.haskell.org/package/dequeue-0.1.12/docs/src/Data-Deque.html#check>